This Instructor's Manual and/or Test Bank is intended only for the course-related work of instructors using Logic and Philosophy: A Modern Introduction, Eleventh Edition. Users are not licensed to place this content on the Internet unless the site is protected by a password available only to students using this book.

Introduction

This instructor's manual is intended to be used with the eleventh edition of *Logic and Philosophy: A Modern Introduction*. It contains a few introductory comments on the contents and use of the text, sample examination questions and answers, and answers to odd numbered exercises (even numbered exercises are answered in the text).

There are four main changes in the text from the tenth edition. In Chapter One, the sections on consistency have been rewritten. In Chapter Five, a section has been added to discuss the suggestion that a rule allowing a tautology to be inserted anytime in a proof could be used as an alternative to conditional proof. In Chapter Twelve, the flow chart governing predicate truth trees has been changed in order to simplify some trees. The possibility of deviating from the order of the flow chart is given as a second method but is not used in the text (for an example of how the two different methods can result in different trees see the answer to exercise 12-1, #5 in the answers to the exercises below). Finally, the material in Part Three of the tenth edition, which included chapters on informal fallacies, inductive logic, axiom systems and alternative logics, has been removed from the printed version of the book, and will be available only online.

The text is intended for use in introductory and intermediate symbolic logic courses. It can be used to support a variety of syllabi, depending on time limitations and the goals of the instructor. For example, it may be used for a standard introductory course covering both traditional and modern deductive logic:

- 1. Sentential Logic Symbolizations and Truth Tables (Chapters 1-3)
- 2. Sentential Logic Proofs (Chapters 4-5)
- 3. Monadic Predicate Logic Symbolizations and Semantics (Chapters 7-8)
- 4. Monadic Predicate Logic Proofs (Chapter 9)
- 5. Syllogistic Logic (Chapter 14)

For a somewhat less demanding course of this type, some instructors may choose to do the chapters on truth trees (Chapters 6 and 12) rather than the chapters on proofs (Chapters 4, 5, 9), or they may choose to omit some of the more philosophical sections at the end of the chapters, such as section 10 of Chapter Three, and sections 6-8 of Chapter Five. Generally, the chapters are constructed so that later sections that involve more philosophical issues (see also, sections 4-5 of Chapter Six, section 10 of Chapter Ten, section 6 of Chapter 12, sections 4-7 of Chapter 13, and section 14 of Chapter 14) or alternative methods (see section 9 of Chapter 5, section 11 of Chapter 10, and section 4 of

Chapter 11) may be omitted if extra time is wanted for covering material in other chapters. Alternatively, the standard course could be made more demanding by including the chapters on trees as well as proofs, or by including material on Relational Predicate Logic from Chapter 10.

Some programs only offer a symbolic logic class at the intermediate level. For such a course the more demanding version of the standard course would be appropriate, as well as one that omitted the material from Chapter 14 and instead covered the material in Chapter 13 on identity and relations:

- 1. Sentential Logic Symbolizations and Truth Tables (Chapters 1-3)
- 2. Sentential Logic Proofs (Chapters 4-5)
- 3. Monadic Predicate Logic Symbolizations and Semantics (Chapters 7-8)
- 4. Monadic Predicate Logic Proofs (Chapter 9)
- 5. Relational Predicate Logic (Chapters 10, 11 & 13)

One may also choose to do the chapters on truth trees instead of some of the more philosophical sections or the sections involving alternative methods.

Incorporation of some of the online material allows for an introductory course which presents the students with the fundamentals of both deductive and inductive logic. Such a course of this type could be structured as follows:

- 1. Sentential Logic Symbolizations and Truth Tables (Chapters 1-3)
- 2. Monadic Predicate Logic Symbolizations and Semantics (Chapters 7-8)
- 3. Syllogistic Logic (Chapter 14)
- 4. Informal Fallacies (Chapter 15 online)
- 5. Inductive Logic (Chapter 16 online)

Such a course could be followed by a second semester course that covered the more demanding material of the text:

- 1. Sentential Logic Proofs and Trees (Chapters 4-6)
- 2. Monadic Predicate Logic Proofs (Chapter 9)
- 3. Relational Predicate Logic Proofs and Predicate Trees (Chapters 10-12)
- 4. Identity and Philosophical Problems of Symbolic Logic (Chapter 13)
- 5. Axiom Systems and Alternative Logics (Chapters 17-18 online)

Of course, many other ways of using the text are possible.

As is clear from the outlines above, traditional syllogistic logic is placed in Chapter 14 after the presentation of modern sentential and predicate logic. Traditional logic obviously came first. Some texts approach the more modern symbolic logic by way of traditional logic. But historical order does not necessarily mirror sound pedagogical order. In fact, most instructors who have tried it both ways have found it best to reverse historical order and start with modern sentential and predicate logic. In particular, it has

been found that students grasp the traditional material much more swiftly if they first have been exposed to modern predicate logic.

Consequently, sentential and predicate logic have been covered in the text before the traditional material. But the section on traditional logic has been written so that it can be covered first with only a very small loss of content. However, instructors are urged to try teaching the material in the order presented in the text if they have not previously tried it this way, in order to experience firsthand the great pedagogical advantage that accrues from doing so.

In learning logic the exercises are especially vital. Experience has shown that it is a rare student indeed who can successfully absorb the material in this text without doing at least a good sample of the exercises. This text has been designed to provide the student with significant help with the exercises by means of walk-through sections preceding some of the most crucial exercise sets.

Finally, students should be urged to use the Key Terms sections at the end of each chapter, first in their original mastery of the material and again for review. Students who truly understand all of the terms covered are well on their way toward understanding all of the material in the book. (Key terms occur in the text in boldface on their first or most important occurrences.)

Sample Examination Questions and Answers

I. CHAPTER ONE: INTRODUCTION

- 1. Consider the following argument:
 - 1. All men are authors.
 - 2. Mark Twain is a man.
 - /:. 3. Mark Twain is an author.
 - a. Are all of the premises true?
 - b. Is the conclusion true or false?
 - c. Is the argument deductively valid or deductively invalid?
 - d. Is the argument sound? Explain, giving the definition of a sound argument.
- 2. Consider the following argument:
 - 1. If George Bush is a Republican, then he is not a Democrat.
 - 2. George Bush is not a Democrat.
 - 3. George Bush is a Republican.
 - a. Are all of the premises true?
 - b. Is the conclusion true or false?
 - c. Is the argument deductively valid or deductively invalid?
 - d. Is the argument sound? Explain, giving the definition of a sound argument.
- 3. Suppose you know of an argument only that it is valid and has a true conclusion. What, if anything, can you tell about its premises? (Defend your answer, including examples.)
- 4. Suppose you know of an argument only that it is valid and has a false conclusion. What, if anything, can you tell about its premises? (Defend your answer.)
- 5. Suppose you know that an argument is sound. What can you determine about its conclusion? (Defend your answer.)
- 6. Suppose you know of an argument only that it has all true premises and a true conclusion. Can you tell from that whether it is valid or invalid? (Defend your answer.)
- 7. Suppose you know that argument is invalid. Can you tell from that whether its conclusion is true or false? (Defend your answer.)
- 8. Suppose you know that an argument is valid. Can you tell from that whether its conclusion is true or false? (Defend your answer.)
- 9. Suppose you know that a set is consistent. Can you tell from that whether every set member is actually true? (Defend your answer.)

10. Suppose you know every member of a set of sentences is false. Can you tell from that whether the set is inconsistent? (Defend your answer.)

I. ANSWERS

- 1. a. No, the first premise is false.
 - b. True.
 - c. Valid.
 - d. No, since a sound argument is an argument that is valid and has all true premises, but this argument has a false premise.
- 2. a. Yes.
 - b. True.
 - c. Invalid, since it is possible that the first two premises are true but the conclusion is false (for example if George Bush were an independent).
 - d. No, since a sound argument is an argument that is valid and has all true premises, but this argument is invalid.
- 3. You can't tell anything about its premises. For instance, the valid argument "If Socrates is a man then Socrates is mortal. Socrates is a man. Therefore, Socrates is mortal", has a true conclusion and all true premises. But the valid argument "If Koko (the gorilla) is a man, then Koko is mortal. Koko is a man. Therefore, Koko is mortal", has a true conclusion, and one true and one false premise.
- 4. At least one of the premises must be false since a valid argument cannot have all true premises and a false conclusion.
- 5. If an argument is sound, then it is valid and has all true premises. If an argument is valid then it is impossible for it to have all true premises and a false conclusion. Therefore, the conclusion of a sound argument must be true.
- 6. No, both valid and invalid arguments can have all true premises and a true conclusion. You would need to determine if it is also possible for the argument to have all true premises and a false conclusion.
- 7. No, the conclusion of an invalid argument could be true or false. We only know that in either case it will be possible for it to have all true premises and a false conclusion.
- 8. No, the conclusion of an valid argument could be true or false. We only know that in either case it will be impossible for it to have all true premises and a false conclusion.
- 9. No, a person can have a consistent set of beliefs where one or more are false. If the set is consistent, it only follows that it is possible for every set member to be true.

10. No. You would need to determine whether it is impossible for all members of the set to be true.

II. CHAPTER TWO: SYMBOLIZING IN SENTENTIAL LOGIC

A. General Theory

For 1-3, circle one of a-d:

- 1. A compound sentence is truth–functional if and only if:
 - a. the truth value of the compound sentence is determined by the truth values of its component sentences
 - b. each of its component sentences is either true or false
 - c. the truth values of the component sentences are determined by the truth values of the compound
 - d. none of these
- 2. The sentence "It will rain only if the temperature drops" is correctly symbolized (using obvious abbreviations) as:
 - a. $R \supset D$
 - b. $D\supset R$
 - c. $R \equiv D$
 - d. none of these
- 3. The sentence "It will rain unless the temperature drops" is correctly symbolized (using obvious abbreviations) as:
 - a. $R \supset D$
 - b. $D \supset R$
 - c. $R \equiv D$
 - d. none of these
- 4. a. Give an original example of a *truth*—*functional* use of a sentence connective in ordinary English, and defend your view that it is truth—functional.
- b. Give an original example of a *non–truth–functional* use of a sentence connective in ordinary English, and defend your view that it is non–truth-functional.

For 5 and 6, identify the main connective.

5.
$$\sim [(A \lor B) \equiv (B \cdot C)] \supset (F \lor G)$$

6.
$$\sim [\sim G \supset (B \lor A)]$$

A. Answers

- 1. a
- 2. a

3. d

4. a. The word "and" in the sentence "Art went to the show and Betsy went to the show" is truth—functional, because the truth value of the sentence formed by this connective is a function of the truth value of the two conjuncts, "Art went to the show" and "Betsy went to the show".

b. The word "because" in the sentence "Art went to the show because Betsy went to the show" is not truth—functional. For example the sentence may be false if both component sentences are true or it may be false. In other words, the truth value of the sentence is not a function of its component parts.

5. The horseshoe.

6. The first tilde.

B. Symbolizing

Symbolize the following, letting R = "It will rain", U = "The temperature will go up", D = "The temperature goes down", S = "It will snow":

1. It won't snow, unless the temperature goes down.

2. The temperature will go up only if it doesn't rain, and will go down if it does rain.

3. The temperature will go down just in case it rains.

4. The temperature will go up or down but not both.

5. If it neither snows nor rains, the temperature won't go up and it won't go down.

6. Either it snows and the temperature won't go up, or it rains and the temperature won't go down.

7. If it rains, then if the temperature goes down it will snow.

8. If it rains, then it snows if and only if the temperature goes down.

9. It won't rain, but it will snow unless the temperature goes up.

10. If it rains, the temperature will go up unless, of course, it goes down.

11. It will snow if and only if it does not rain and the temperature does not go up.

12. If it both rains and snows, then the temperature is neither going up nor going down.

B. Answers

1. $\sim D \supset \sim S$

2. $(U \supset \sim R) \cdot (R \supset D)$

3. $D \equiv R$

4. $(U \lor D) \cdot \sim (U \cdot D)$

5. $\sim (S \vee R) \supset (\sim U \cdot \sim D)$

6. $(S \cdot \sim U) \vee (R \cdot \sim D)$

7. $R\supset (D\supset S)$

8. $R \supset (S \equiv D)$

9. $\sim R \cdot (\sim U \supset S)$

10. $R \supset (\sim D \supset U)$

11. $S \equiv (\sim R \cdot \sim U)$

12. $(R \cdot S) \supset \sim (U \vee D)$

C. Translating

Translate the following into more or less colloquial English sentences. Let J = "Art gets a new job"; B = "Art gets a new boss"; A = "Art gets a new apartment"; R = "Art get a new roommate."

1.
$$\sim (J \vee B)$$

4.
$$\sim J \supset \sim [(B \lor (A \lor R))]$$

2.
$$J \supset (B \cdot A)$$

5.
$$(\sim A \lor \sim R) \supset \sim J$$

3.
$$\sim (A \cdot R)$$

6.
$$(B \equiv J) \cdot (R \equiv A)$$

C. Answers

- 1. Art will get neither a new job nor a new boss.
- 2. If Art gets a new job then he will get a new boss and a new apartment.
- 3. Art will not get both a new apartment and a new roommate.
- 4. If Art does not get a new job, then he will get neither a new boss, nor a new apartment, nor a new roommate.
- 5. If Art either does not get a new apartment or does not get a new roommate, then he does not get a new job.
- 6. Art gets a new boss if and only if he gets a new job, and a new roommate just in case he gets a new apartment.

III. CHAPTER THREE: TRUTH TABLES

A. General Theory

1. Determine the sentence forms of which the following are substitution instances:

a.
$$\sim (A \equiv B) \supset C$$

b.
$$A \lor (B \cdot \sim C)$$

- 2. If a sentence form contains five variables, how many lines or rows must its complete truth table analysis have?
- 3. Assume you know of an argument only that its premises are *not* consistent. What, if anything, can you tell about the argument's validity? (Defend your answer.)
- 4. Why is it true that any argument with a conclusion with the form $p \lor \sim p$ is valid?
- 5. a. We can define a tautologous sentence as one that is a substitution instance of some tautologous sentence form, and a contradictory sentence as one that is a substitution instance of some contradictory sentence form. Why can't we analogously define a contingent sentence as one that is a substitution instance of some contingent sentence form? (Defend your answer, including examples.)
 b. We cannot define a contingent sentence as one that is a substitution instance of some contingent sentence form. But then how *can* we define that term?

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6. Suppose one of the premises of an argument is logic equivalent to the conclusion. What, if anything, can you conclude about the argument's validity?

A. Answers

1. a.
$$p, p \supset q, \sim p \supset q, \sim (p \equiv q) \supset r$$

b. $p, p \lor q, p \lor (q \cdot r), p \lor (q \cdot \sim r)$

- 2. There would be $2^5 = 32$ rows of the truth table.
- 3. If an argument's premises are not consistent then it will be impossible for all of them to be true. Therefore, the argument must be valid, because it will be impossible for the premises to be true and the conclusion false.
- 4. Any argument with a conclusion of the form $p \lor \sim p$ will have a conclusion that is a tautology, and thus a conclusion that cannot be false. Thus, it will be impossible for that argument to have all true premises and a false conclusion, and so it will be valid.
- 5. a. Because *every* sentence, tautologous, contradictory, or contingent, is a substitution instance of some contingent sentence form or other. For instance, every sentence is a substitution instance of the contingent sentence form *p*.b. We can define a contingent sentence as a sentence that is not a substitution instance of any tautological or contradictory sentence form.
- 6. If one of the premises of an argument is logically equivalent to the conclusion, then the premise and the conclusion will always have the same truth value. Thus, it will be impossible for that argument to have all true premises and a false conclusion, and so it will be valid.

B. Tautologies, Contradictions, and Contingent Sentences

Determine by truth table analysis which of the following sentence forms are tautologous, which are contradictory, and which are contingent:

1.
$$(p \lor q) \equiv [(p \cdot \sim q) \lor q]$$

4.
$$(p \supset \sim q) \supset \sim (q \supset \sim p)$$

2.
$$(p \supset q) \lor (q \supset r)$$

5.
$$\sim [(\sim p \lor q) \equiv (\sim q \supset \sim p)]$$

3.
$$p \equiv (q \equiv p)$$

B. Answers

1. Tautologous

4. Contingent

2. Tautologous

5. Contradictory

3. Contingent

C. Logical Equivalences

Use a truth table to determine which of the following pairs of sentence forms are logically equivalent.

1.
$$\sim (p \supset q), p \cdot \sim q$$

2.
$$\sim (p \equiv q), (\sim p \equiv q) \vee (p \equiv \sim q)$$

3.
$$p \supset (q \cdot r), (p \supset q) \cdot (p \supset r)$$

4.
$$\sim [(p \vee q) \supset p], \sim [(p \supset q) \vee q]$$

C. Answers

- 1. Equivalent
- 2. Equivalent

- 3. Equivalent
- 4. Not Equivalent

D. Proving Validity of Argument Forms

Determine by truth table analysis which of the following argument forms are valid and which are invalid:

$$1.(\sim p \lor q) \lor (\sim p \cdot q)$$

$$2.q \supset p/... \sim q$$

$$1.(p\cdot q)\equiv (q\cdot r)$$

$$2.\sim (r\supset \sim q)/...\sim r\supset p$$

$$1.p \supset \sim (\sim q \lor r)$$

$$2.\sim r \supset q/...\sim p \cdot q$$

1.
$$\sim (p \cdot \sim q) \equiv \sim p$$

$$/: (\sim q \vee p) \vee (\sim q \supset \sim p)$$

D. Answers

- (1) Invalid
- (2) Invalid

- (3) Valid
- (4) Valid

E. Short Truth Table Test for Invalidity

Use the short truth table method to show that the following arguments are invalid:

$$1.A \supset B$$

$$2.C \supset \sim B$$

$$3 \sim C / :: A$$

$$1.A \lor (B \cdot \sim C)$$

$$2.\sim [A\cdot\sim (C\vee B)]$$

$$3.B \supset C$$

$$4.\sim B/:.\sim A$$

$$1.A \supset (B \lor C)$$

$$2.B \supset (C \supset D)$$

$$3.\sim D/:.\sim A$$

$$1.\sim (A\vee B)\supset \sim C$$

 $2.\sim (\sim A \cdot \sim C) \vee \sim B$

 $3.\sim (B\vee D)$

E. Answers

(1)
$$V(A) = \mathbf{F}, V(B) = \mathbf{T}, V(C) = \mathbf{F}$$

 $V(A) = \mathbf{F}, V(B) = \mathbf{F}, V(C) = \mathbf{F}$

(2)
$$V(A) = T$$
, $V(B) = T$, $V(C) = F$, $V(D) = F$
 $V(A) = T$, $V(B) = F$, $V(C) = T$, $V(D) = F$

(3)
$$V(A) = T$$
, $V(B) = F$, $V(C) = T$

(4)
$$V(A) = F$$
, $V(B) = F$, $V(C) = F$, $V(D) = F$, $V(E) = T$

F. Short Truth Table Test for Consistency

Use the short truth table method to show that the following sets of premises are consistent:

(1) 1.
$$(A \lor B) \cdot \sim C$$

2.
$$C \equiv D$$

3.
$$D \lor (A \equiv B) / \therefore B$$

(3) 1.
$$A \cdot (B \equiv E)$$

2.
$$(C \lor D) \supset {}^{\sim} A$$

3.
$$B \supset (C \cdot D) / :: E$$

 $4.B \vee (E \cdot \sim A)$

 $/ :: \sim A \supset C$

(2) 1.
$$D \equiv (A \lor C)$$

2.
$$\sim [B \supset (D \supset E)]$$

3.
$$\sim B \vee \sim C / :. \sim D$$

(4) 1.
$$\sim B \vee C$$

2.
$$(D\supset C)\cdot (A\supset \sim C)$$

3.
$$(C \cdot \sim C) \vee B$$

4.
$$B \supset (D \equiv A) / :. B$$

F. Answers

(1)
$$V(A) = T$$
, $V(B) = T$, $V(C) = F$, $V(D) = F$

(2)
$$V(A) = T$$
, $V(B) = T$, $V(C) = F$, $V(D) = T$, $V(E) = F$

(3)
$$V(A) = T$$
, $V(B) = F$, $V(C) = F$, $V(D) = F$, $V(E) = F$

(4)
$$V(A) = F$$
, $V(B) = T$, $V(C) = T$, $V(D) = F$

IV. CHAPTER FOUR: PROOFS WITHOUT CP OR IP

Prove valid, using the eighteen valid argument forms (but *not* **CP** or **IP**):

(1) 1.
$$A \supset \sim A / :: \sim A$$

$$1.\sim A \lor B$$
$$2.C \supset A$$
$$3.\sim B / \therefore \sim C$$

(3) 1.
$$A/::B\supset (\sim A\supset C)$$

$$1.(A \lor B) \supset C$$
$$2.\sim A \cdot \sim C / \therefore \sim B$$

1.~
$$[A \cdot \sim (\sim A \lor \sim B)]$$

2.~ $(A \cdot B) \supset (C \cdot D) / \therefore D$

1.
$$(A \lor B) \supset C$$

2. $\sim C$
3. $\sim D \lor B$
4. $\sim A \supset \sim E / \therefore \sim (\sim E \supset D)$

$$1.(A \lor B) \supset (C \cdot D) / \therefore B \supset C$$

$$1.(A \cdot B) \supset C$$
$$2.(B \cdot C) \supset D$$
$$3. \sim (E \lor \sim B) / \therefore \sim (\sim D \cdot A)$$

1.
$$A \supset B$$

2. $C \supset D / : (A \cdot C) \supset (B \cdot D)$

IV. ANSWERS

- (l) 1. $A \supset \sim A / \therefore \sim A$ p 2. $\sim A \lor \sim A$ 1 Impl 3. $\sim A$ 2 Taut
- (3) 1. $A/: B \supset (\sim A \supset C)$ p2. $A \lor (B \supset C)$ 1 Add 3. $\sim \sim A \lor (B \supset C)$ 2 DN 4. $\sim A \supset (B \supset C)$ 3 Impl 5. $(\sim A \cdot B) \supset C$ 4 Exp 6. $(B \cdot \sim A) \supset C$ 5 Comm 7. $B \supset (\sim A \supset C)$ 6 Exp

- (5) 1. $\sim [A \cdot \sim (\sim A \vee \sim B)]$ p 2. $\sim (A \cdot B) \supset (C \cdot D) / :: D$ p 3. $\sim [A \cdot (\sim \sim A \cdot \sim \sim B)]$ $1~{\bf DeM}$ 4. $\sim [A \cdot (A \cdot B)]$ 3 **DN** (2x) 5. $\sim [(A \cdot A) \cdot B]$ 4 Assoc 6. $\sim (A \cdot B)$ 5 Taut 7. $C \cdot D$ 2,6 **MP** 8. D 7 Simp
- (6) 1. $(A \lor B) \supset C$ p 2. ~ *C* p 3. $\sim D \vee B$ p 4. $\sim A \supset \sim E / : \sim (\sim E \supset D)$ p 5. $\sim (A \vee B)$ 1,2 **MT** 6. $\sim A \cdot \sim B$ **5 DeM** 7. $\sim A$ 6 Simp 8. $\sim E$ 4,7 **MP** 9. ∼ *B* 6 Simp 10. ∼*D* 3,9 **DS** 11. $\sim E \cdot \sim D$ 8,10 **Conj** 12. $\sim (E \vee D)$ 11 **DeM** 13. \sim (\sim \sim $E \vee D$) 12 **DN** 14. \sim ($\sim E \supset D$) 13 **Impl**
- **(7)** 1. $(A \lor B) \supset (C \cdot D) / \therefore B \supset C$ p 2. $\sim (A \vee B) \vee (C \cdot D)$ 1 Impl 3. $(\sim A \cdot \sim B) \vee (C \cdot D)$ 2 **DeM** 4. $[(\sim A \cdot \sim B) \lor C] \cdot [(\sim A \cdot \sim B) \lor D]$ 3 **Dist** 5. $(\sim A \cdot \sim B) \vee C$ 4 Simp 6. $C \vee (\sim A \cdot \sim B)$ 5 Comm 7. $(C \lor \sim A) \cdot (C \lor \sim B)$ 6 Dist 8. $C \lor \sim B$ 7 Simp 9. $\sim B \vee C$ 8 Comm 10. $B \supset C$ 9 **Impl**

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(8)
       1. (A \cdot B) \supset C
                                                        p
        2. (B \cdot C) \supset D
                                                         p
        3. \sim (E \vee \sim B) / \therefore \sim (\sim D \cdot A)
                                                         p
        4. \sim E \cdot \sim \sim B
                                                         3 DeM
        5. \sim E \cdot B
                                                         4 DN
        6. B
                                                         5 Simp
        7. B\supset (C\supset D)
                                                         2 Exp
        8. C \supset D
                                                         6, 7 MP
        9. (A \cdot B) \supset D
                                                         1, 8 HS
       10. (B \cdot A) \supset D
                                                         9 Comm
       11. B\supset (A\supset D)
                                                         10 Exp
       12. A \supset D
                                                         6, 11 MP
       13. \sim A \vee D
                                                         12 Impl
       14. \sim A \lor \sim \sim D
                                                         13 DN
       15. \sim (A \cdot \sim D)
                                                         14 DeM
      16. \sim (\sim D \cdot A)
                                                         15 Comm
(9)
        1. A \supset B
                                                                      p
        2. C \supset D / :: (A \cdot C) \supset (B \cdot D)
                                                                      р
                                                                      1 Impl
        3. \sim A \vee B
        4. (\sim A \vee B) \vee \sim C
                                                                      3 Add
        5. \sim A \vee (B \vee \sim C)
                                                                      4 Assoc
        6. \sim A \vee (\sim C \vee B)
                                                                      5 Comm
        7. (\sim A \lor \sim C) \lor B
                                                                      6 Assoc
        8. \sim C \vee D
                                                                      2 Impl
        9. (\sim C \lor D) \lor \sim A
                                                                      8 Add
       10 \sim A \vee (\sim C \vee D)
                                                                      9 Comm
       11. (\sim A \lor \sim C) \lor D
                                                                     10 Assoc
       12. [(\sim A \lor \sim C) \lor B] \cdot [(\sim A \lor \sim C) \lor D] 7,11 Conj
      13. (\sim A \lor \sim C) \lor (B \cdot D)
                                                                     12 Dist
       14. \sim (A \cdot C) \vee (B \cdot D)
                                                                     13 DeM
       15. (A \cdot C) \supset (B \cdot D)
                                                                     14 Impl
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V. CHAPTER FIVE: PROOFS WITH **CP** OR **IP**

A. General Theory

1. Suppose you know that a particular two–premise argument is invalid. Now suppose we add the negation of the conclusion of the two premises to form a three–sentence set of premises. Can a contradiction be derived from this three–sentence set of premises? (Defend your answer.)

2. a. Use **IP** to prove that the following argument is valid.

$$A \supset B$$

 $A \supset \sim B / : \sim A$

b. To illustrate how indirect proofs are a kind of shortened conditional proof, cross out the last line in the above proof and complete it as a conditional proof. (Hint: as an intermediate step prove $A \supset \sim A$.)

A. Answers

 No, because derivation of a contradiction would constitute an *indirect proof* of validity for the argument, but by the hypothesis of the problem, the argument in question is invalid.

2. a. 1.
$$A \supset B$$

2. $A \supset \sim B$
3. A AP/:. $\sim A$
4. B 1,4 MP
5. $\sim B$ 2,4 MP
6. $B \cdot \sim B$ 5,6 Conj
7. $\sim A$ 3-7 IP
b. 7. $B \lor \sim A$ 4 Add
8. $\sim A$ 5,7 DS
9. $A \supset \sim A$ 3-8 CP
10. $\sim A \lor \sim A$ 9 Impl
11. $\sim A$ 10 Taut

B. Proofs with **CP** or **IP**

Prove valid, using the eighteen valid argument forms and **CP** or **IP**:

$$\begin{aligned}
1.A \supset B & 1.A \lor B \\
2.C \supset D & 2.C \supset \sim A \\
/ \therefore (A \cdot C) \supset (B \cdot D) & 3.D \supset E \\
4. \sim D \supset C \\
1. (A \cdot B) \supset C & 5.E \supset \sim A / \therefore B \\
2. (A \cdot \sim B) \supset \sim C & 1. A \supset (B \supset C) \\
/ \therefore A \supset (B \equiv C) & 1. A \supset (B \supset C) \\
2. \sim C \supset (A \cdot B) / \therefore C
\end{aligned}$$

$$\begin{array}{ll}
1.(A \lor B) \supset & 1. \sim (A \cdot \sim B) \\
[(C \lor D) \supset E] & 2. \sim [\sim C \cdot (\sim A \cdot \sim D)] \\
/ \therefore A \supset & 3. \sim [A \cdot (B \cdot \sim D)] \\
[\sim E \supset \sim (C \cdot D)] & / \therefore D \lor C
\end{array}$$

B. Answers