TRUE/FALSE

1.	Early memory management schemes are still used in today's operating systems.					
	ANS: F PTS: 1 REF: 32					
2.	All computers have only a finite amount of memory and if a program doesn't fit, then either the size of the main memory must be increased or the program must be modified.					
	ANS: T PTS: 1 REF: 32					
3.	To overlay is to transfer segments of a program from main memory into secondary storage for execution, so that two or more segments take turns occupying the same memory locations.					
	ANS: F PTS: 1 REF: 32					
4.	The first step in loading a job in a single-user system is storing the first memory location of program into the base register (for memory protection).					
	ANS: T PTS: 1 REF: 33					
5.	A single-user system supports multiprogramming.					
	ANS: F PTS: 1 REF: 33					
6.	The first attempt to allow for multiprogramming used fixed partitions.					
	ANS: T PTS: 1 REF: 34					
7.	The problem of partition intrusion is present in single-user contiguous allocation schemes.					
	ANS: F PTS: 1 REF: 34					
8.	The algorithm used to store jobs into memory requires a few more steps than the one used for a single-user system because the size of the job must be matched with the size of the partition to make sure it fits completely.					
	ANS: T PTS: 1 REF: 34					
9.	The fixed partition scheme does not require that the entire program be stored contiguously and in memory from the beginning to the end of its execution.					
	ANS: F PTS: 1 REF: 35					
10.	. The fixed partition scheme works well if all of the jobs run on the system are of the same size or if the sizes are known ahead of time and don't vary between reconfigurations.					
	ANS: T PTS: 1 REF: 35					

11.	In a fixed partition scheme, large jobs may have a longer turnaround time as they wait for free partitions of sufficient size or may never run.				
	ANS: T	PTS:	1	REF:	36
12.	The best-fit allocation memory to high-order		•	ee/busy	lists organized by memory locations, low-order
	ANS: F	PTS:	1	REF:	38
13.	A large job can have	probler	ns with a first-	fit mem	ory allocation list.
	ANS: T	PTS:	1	REF:	39
14.	The best-fit free list to implement.	scheme	uses memory r	nore ef	ficiently than the first-fit free scheme but it is slower
	ANS: T	PTS:	1	REF:	40
15.	The first-fit algorithm blocks.	n assun	nes that the Me	mory M	Sanager keeps only one list containing free memory
	ANS: F	PTS:	1	REF:	40
16.	6. One of the problems with the best-fit algorithm is that the entire table must be searched before the allocation can be made because the memory blocks are physically stored in sequence according to thei location in memory.				
	ANS: T	PTS:	1	REF:	42
17.	Research continues t	o focus	on finding the	optimu	m allocation scheme.
	ANS: T	PTS:	1	REF:	43
18.	For a fixed partition	system,	memory deallo	ocation	is quite complex.
	ANS: F	PTS:	1	REF:	44
19.	A null entry in the bureturned to the free l		occurs when a 1	memory	y block between two other busy memory blocks is
	ANS: T	PTS:	1	REF:	47
20.		npty blo	ocks and compa	ct them	Memory Manager relocates programs to gather to make one block of memory large enough to t in.
	ANS: T	PTS:	1	REF:	48
21.	Memory is allocated	during	garbage collect	tion.	
	ANS: F	PTS:	1	REF:	48 49

22.	During compaction, the operating system must distinguish between addresses and data values, and the distinctions are not obvious once the program has been loaded into memory.					
	ANS: T	PTS: 1	REF:	49		
23.	After relocation and	compaction,	both the free list	and the busy list are updated.		
	ANS: T	PTS: 1	REF:	51		
24.	The bounds register in memory accessibl		- '	lowest, depending on the specific system) location		
	ANS: T	PTS: 1	REF:	52		
25.	Compaction should a	always be pe	erformed only whe	n there are jobs waiting to get in.		
	ANS: F	PTS: 1	REF:	53 54		
MUL	ГІРЬЕ СНОІСЕ					
1.	 The following,, describes the first memory allocation scheme. a. Each program to be processed was loaded into secondary storage, then swapped into memory in parts b. Each program to be processed was partially loaded into memory, then granted more memory as needed c. Each program to be processed was allocated a portion of memory and could negotiate with other programs to access more memory d. Each program to be processed was loaded in its entirety into memory and allocated as much contiguous space in memory as it needed 					
	ANS: D	PTS: 1	REF:	32		
2.	In a single-user syste a. sequentially b. intermittently		c. d.			
	ANS: A	PTS: 1	REF:	32		
3.	 In the algorithm to load a job in a single-user system, the program counter is initially set to a. the address of the last memory location b. the number of instructions c. zero d. the address of the first memory location 					
	ANS: D	PTS: 1	REF:	33		
4.	Fixed partitions are a a. complete b. static	also called _	partitions. c. d.			
	ANS: B	PTS: 1	REF:	34		
5.	is the first step a. Comparing job s		thm to load a job flargest partition	in a fixed partition.		

	c.	Determining the Setting counter the Placing the job is	o one	-	y size	
	ANS	S: B	PTS:	1	REF:	34
6.	the	·				Manager uses to keep track of jobs is composed of
	b. c.	partition size, me status, access, ar partition size, sta partition size, me	nd memo atus, and	ory address I access		tus
	ANS	S: D	PTS:	1	REF:	35
7.	a. b. c.	fixed partition s when jobs have when jobs have when job sizes a when all jobs are	the same different re not k	e size t sizes nown in advanc		
	ANS	S: A	PTS:	1	REF:	35
8.		_ consists of frag An inefficient fit Indirect partition	t	f free memory	c.	n blocks of allocated memory. External fragmentation Internal fragmentation
	ANS	S: C	PTS:	1	REF:	36
9.	men a.	keeps the inory. fixed partition al first-fit memory	location	1	c.	dynamic fit memory allocation best-fit memory allocation
		S: B	PTS:		REF:	·
10.		has the least was Fixed partitioning.	ıg		c.	nrtition fitting the requirements. Dynamic fit memory allocation Best-fit memory allocation
	ANS	S: D	PTS:	1	REF:	38
11.	sche Jobs J1 1 J2 2 J3 3 J4 1	eme is used, the j s: 0K 0K 0K 0K				s 1-4 and memory blocks. Assuming a first-fit
	Bloc B1 3 B2 3 B3 3 B4 2 a.	30K 15K 50K			c.	Ј3

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b. J2
                                                  d. J4
                                              REF: 39|40
    ANS: C
                         PTS: 1
12. Consider the following space requirements for jobs 1-4 and memory blocks. Assuming a best-fit
    scheme is used, the job, , is placed in the last block.
    Jobs:
    J1 10K
    J2 20K
    J3 30K
    J4 10K
    Blocks:
    B1 30K
    B2 15K
    B3 50K
    B4 20K
    a. J1
                                                  c. J3
                                                  d. J4
    b. J2
    ANS: B
                         PTS: 1
                                              REF: 40
13. The following algorithm can be described as _____.
     1 Set counter to 1
    2 Do while counter <= number of blocks in memory
      If job size > memory size(counter)
        Then counter = counter + 1
      Else
        load job into memory size(counter)
        adjust free/busy memory lists
        go to step 4
    End do
    3 Put job in waiting queue
    4 Go fetch next job
    a. first-fit memory allocation
                                                  c. least-fit memory allocation
    b. best-fit memory allocation
                                                  d. fixed partition memory allocation
    ANS: A
                         PTS: 1
                                              REF: 41
14. The following algorithm can be described as _____.
     1 Initialize memory_block(0) = 99999
    2 Compute initial memory waste = memory block(0) - job size
    3 Initialize subscript = 0
    4 Set counter to 1
    5 Do while counter <= number of blocks in memory
         If job size > memory size(counter)
            Then counter = counter + 1
         Else
            memory waste = memory size(counter) - job size
         If initial memory waste > memory waste
            Then subscript = counter
            initial memory waste = memory waste
            counter = counter + 1
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	Else Load job into	y allocation	c.	least-fit memory allocation fixed partition memory allocation	
	ANS: B	PTS: 1	REF:	42	
15.	used, is the beginning Address 4075 5225 6785 7560 7600 10250 15125 24500 a. 6785 b. 7600 ANS: B	beginning address Memory	s of the block g Block Size 105 5 600 20 205 4050 230 1000 c. d.	10250 15125 43	
16.	a. Memory Manage. Memory Manage. Memory Manage. Memory Manage. Memory Manage. "free."	ger immediately gger adds block to ger resets the state	lock and comb gives memory free list and re us of the memo	bines it with another free block. to another program. removes it from busy list. nory block where the job was stored to	
	ANS: D	PTS: 1	REF:	44	
17.	two other free blocka. The sizes of theb. All three are moneyc. The block is contact.	ks? e three free partiti oved individually	ons must be co from the busy larger of the tw	Memory Manager deallocates a block that is between combined. by list to the free list. wo adjacent blocks.	en
	ANS: A	PTS: 1	REF:	46	
18.	When memory is do a. blank line b. null entry ANS: B	eallocated, an ent	c.	noved from the free list by creating a(n) joined entry empty entry 46	
19.	A(n) in the burreturned to the free		en a memory b	block between two other busy memory blocks is	
	a. blank line		c.	joined entry	

	b. null entry			d.	empty entry
	ANS: B	PTS:	1 RI	EF:	47
20.				ctions	page collection or defragmentation, is performed by s of the memory space. compaction reallocation
	ANS: C	PTS:	1 RI	EF:	48
21.	The operating system operation code.	n can tel	the of each	grou	up of digits by its location in the line and the
	a. functionb. value				order assignment
	ANS: A	PTS:	1 RI	EF:	49
22.	The contains the able to access the contains the contains the able to access the contains the contains the able to access the ac	rrect me		ter re	each address referenced in the program so it will be elocation. relocation register bounds register
	ANS: C	PTS:	1 RI	EF:	52
23.	is the actual n a. 1,800 b. 18,000	nemory a	address for a job t	c.	tarts at 18K. 18,432 180,000
	ANS: C	PTS:	1 RI	EF:	53
24.	By compacting and r throughput, but an ur a. null entries b. segmentation			ore _ c.	er optimizes the use of memory and thus improves errors overhead
	ANS: D	PTS:	1 RI	EF:	53
25.	One approach to perfa. byte b. percentage	forming	compaction is to o	c.	when a certain of memory becomes busy. bit area
	ANS: B	PTS:	1 RI	EF:	53